

# On characterizing classical and quantum entropy

Arthur Parzygnat IHÉS

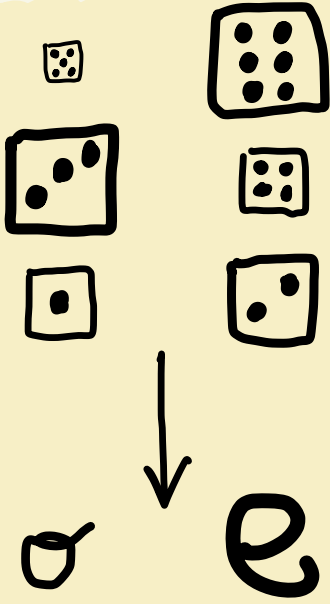
Categorical semantics of Entropy

CUNY Graduate Center ITS

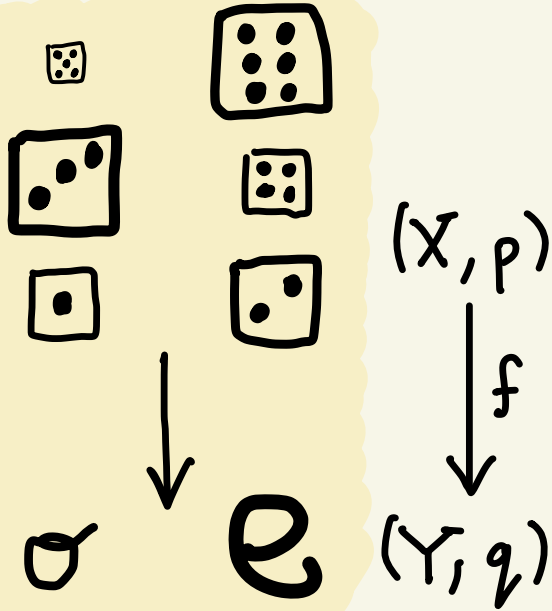
5/13/2022

# Information Loss

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The entropy difference

$$\Delta_H(f) := H(p) - H(q), \text{ where}$$

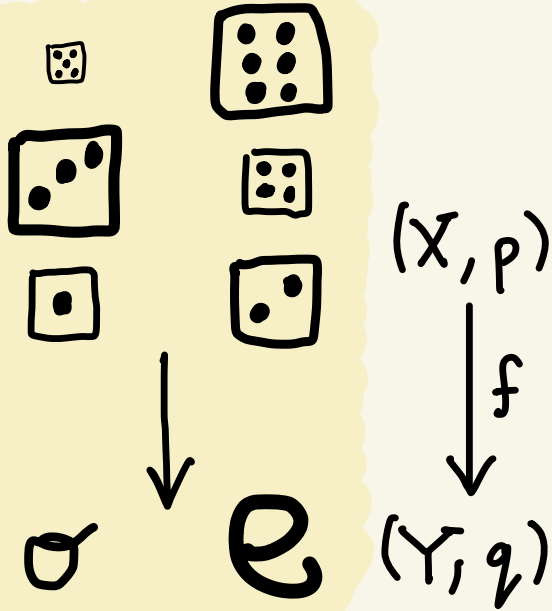
$$H(p) := -\sum_x p_x \log(p_x)$$

is the Shannon entropy, quantifies this.

Here,  $q_y = \sum_{x \in f^{-1}(\{y\})} p_x$  is

the pushforward of  $p$  along  $f$ , denoted  $f_*p$ .

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Baez, Fritz, Leinster characterized  $\Delta_H$

as a continuous convex functor

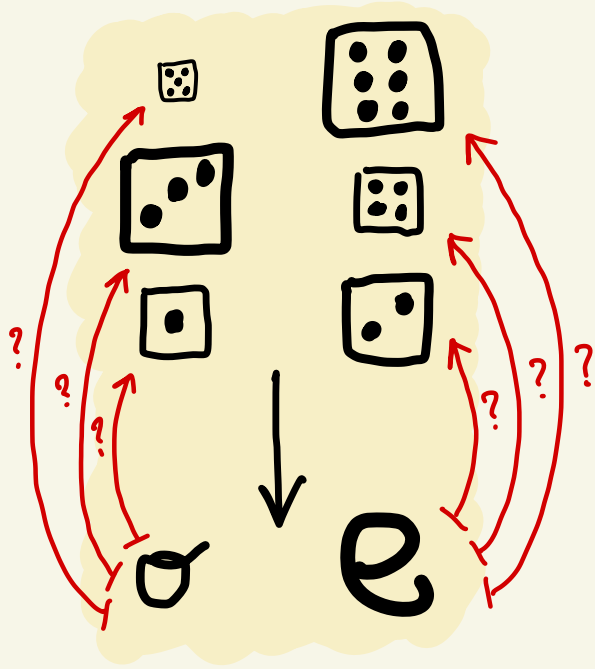
into  $\mathbf{BR}_{\geq 0}$  ( $\mathbb{R}_{\geq 0}$  viewed as a

one-object category) up to a constant  $\geq 0$ .

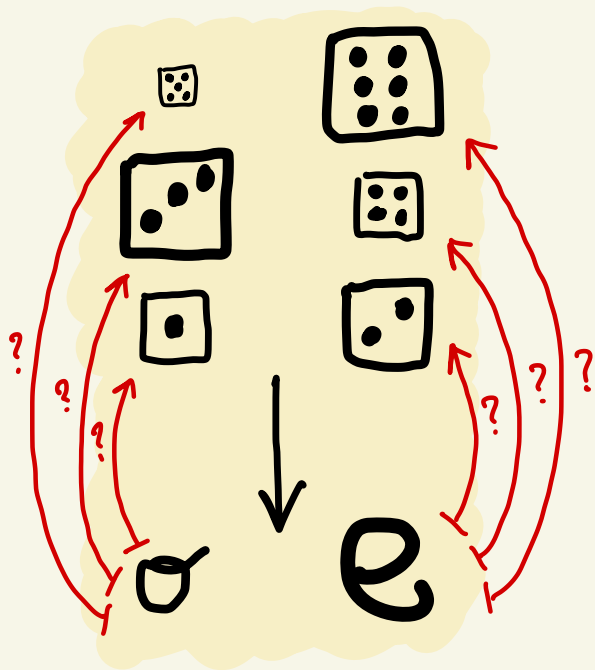


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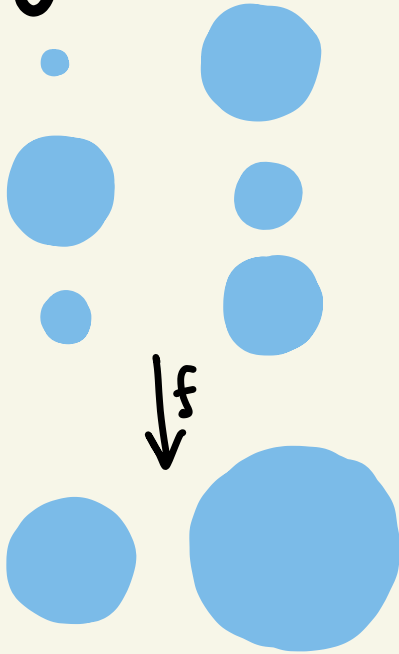
But we can try using a stochastic procedure  $Y \xrightarrow{h} X$ .

## Notation

$h$  assigns a probability distribution  $h_y$  on  $X$  to each  $y \in Y$ .  
The value on  $x \in X$  is written  $h_{xy}$ .

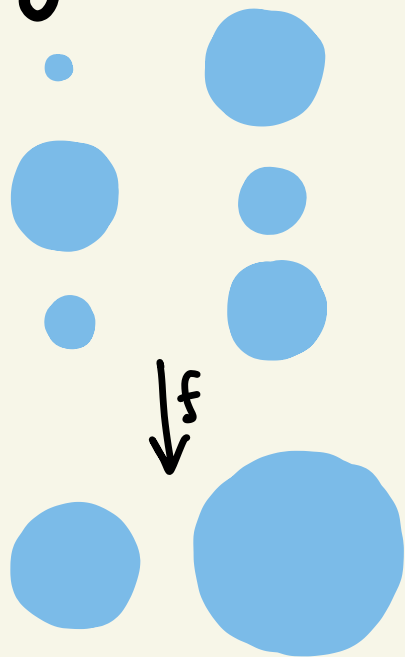
Compare to  $P(x|y)$  the conditional probability of  $x$  given  $y$ .

Hypotheses = stochastic sections

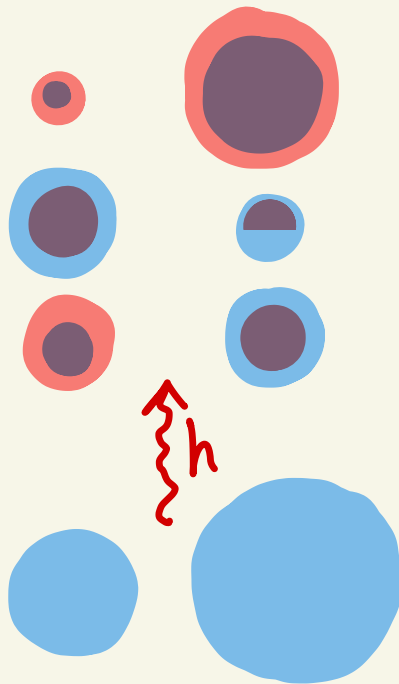


Water droplet picture  
on left due to Gromov

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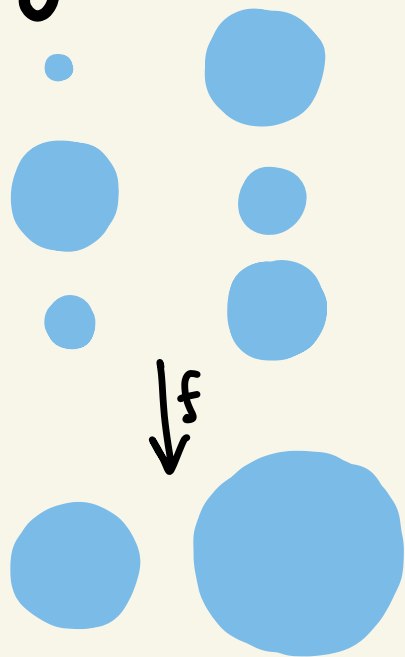


make  
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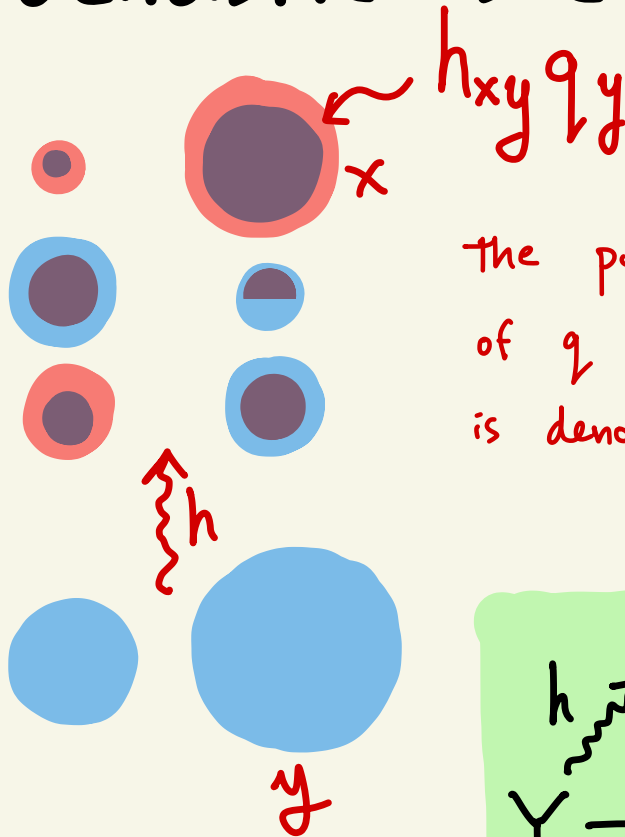
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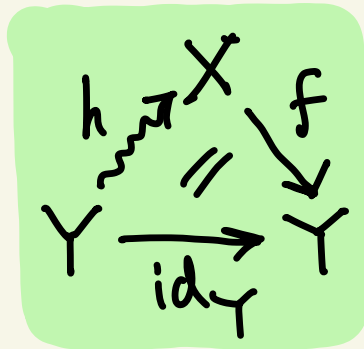
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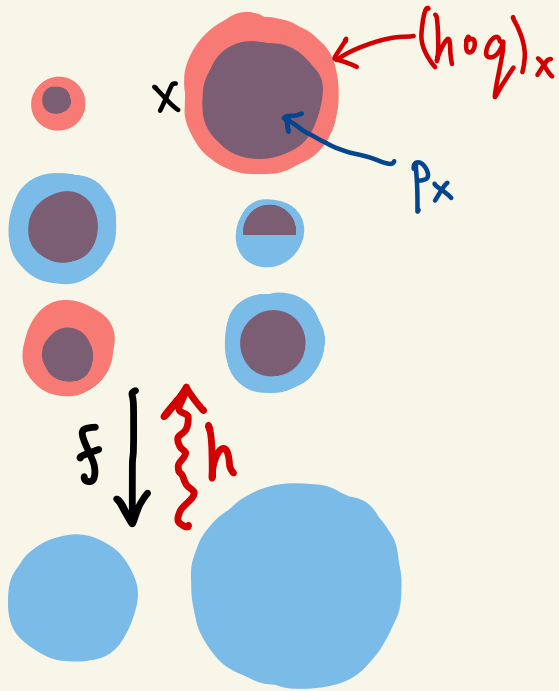


$h_{xy} q_y$

the pushforward  
of  $q$  along  $h$   
is denoted  $h_0 q$ .

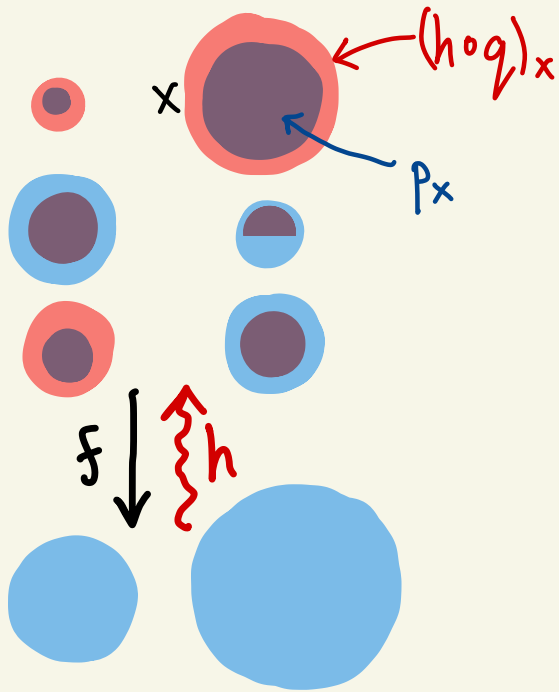


# Relative Entropy of Recovery



But hypotheses are not always correct.

# Relative Entropy of Recovery

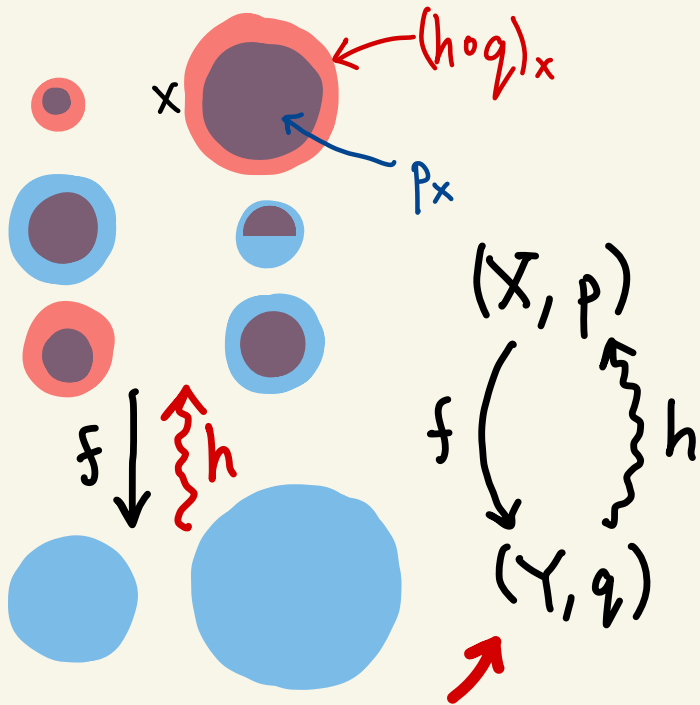


But hypotheses are not always correct. To quantify the accuracy of a hypothesis  $h$ , we use the relative entropy (of recovery)

$$S(p \parallel h \circ q) := \sum_x P_x \log \left( \frac{P_x}{(h \circ q)_x} \right) \\ = \sum_x P_x \log \left( \frac{P_x}{h_{x f(x)} q_{f(x)}} \right)$$

$h$  is sometimes called a recovery map

# Relative Entropy of Recovery



Warning: This notation is misleading.  
 We require  $q = f \circ p$  but NOT  $p = h \circ q$ !

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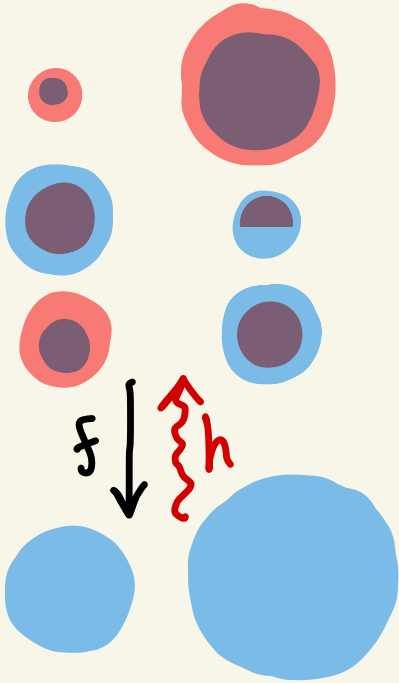
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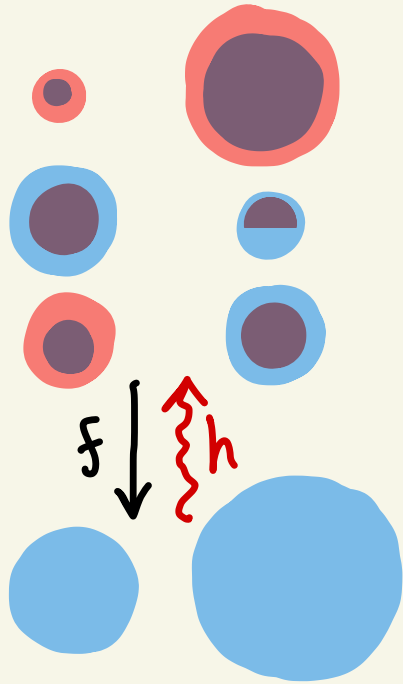
# Relative Entropy & Optimal Hypotheses

The relative entropy satisfies

- $S(p \parallel h \circ g) \geq 0$



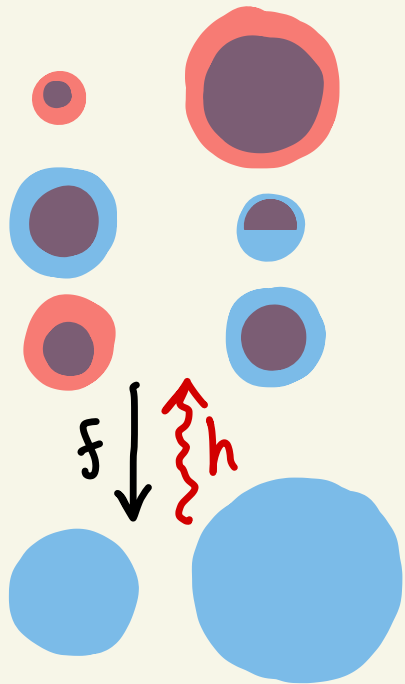
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An optimal hypothesis is a hypothesis  $h$  such that  $h \circ q = p$ .



# A Category of Hypotheses

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morphisms:  $(X, p) \begin{array}{c} \xrightarrow{f} \\ \text{~~~~~} \\ \xleftarrow{h} \end{array} (Y, q)$

$f = \text{deterministic}$

$q = f \circ p$

$f \circ h = \text{id}_Y$

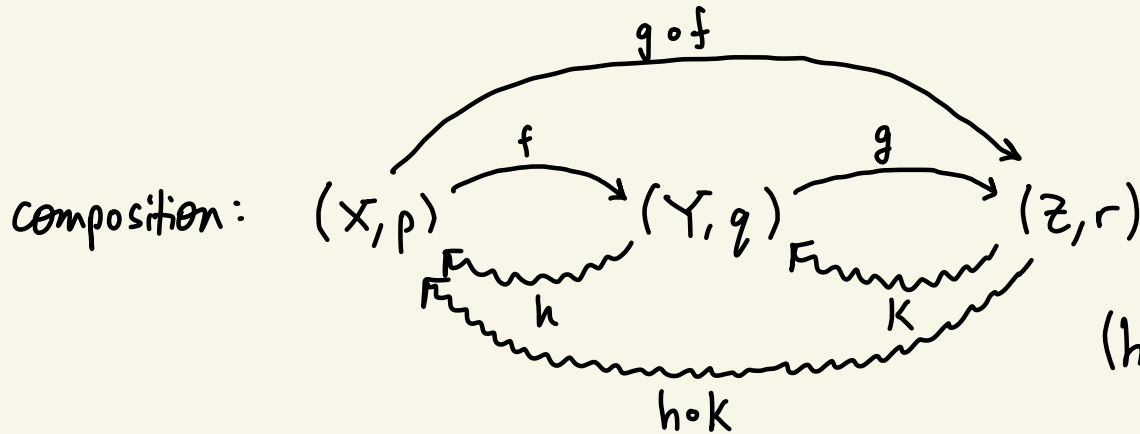
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 $q = f \circ p$   
 $f \circ h = \text{id}_Y$  (BUT NOT  $p = h \circ q$ )



$$(h \circ k)_{xz} := \sum_y h_{xy} k_{yz}$$

Chapman-Kolmogorov



# Relative Entropy as a Functor

$$\text{FinStat} \xrightarrow{\text{RE}} \mathbb{B}[0, \infty]$$

$\mathbb{B}[0, \infty]$  is the category with a single object and the morphism set is  $[0, \infty]$  with addition as the composition.

# Relative Entropy as a Functor

$$\text{FinStat} \xrightarrow{\text{RE}} \mathbb{B}[0, \infty]$$

A commutative diagram illustrating the relationship between probability distributions and relative entropy. On the left, a square diagram shows a mapping  $f$  from  $(X, p)$  to  $(Y, q)$ . A wavy arrow labeled  $h$  points from  $(Y, q)$  back to  $(X, p)$ . An arrow labeled  $S(p \parallel h \circ q)$  points from the square to the right. Above the square, the text  $\text{FinStat} \xrightarrow{\text{RE}} \mathbb{B}[0, \infty]$  indicates the functorial mapping.

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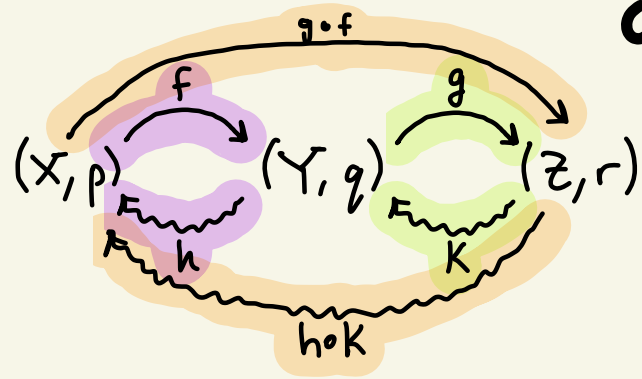
$$\text{FinStat} \xrightarrow{\text{RE}} \mathbb{B}[0, \infty]$$
$$\begin{array}{ccc} (X, p) & & \\ \downarrow f & \rightsquigarrow h & \\ (Y, q) & & \end{array} \quad \longmapsto \quad S(p \parallel h \circ q)$$

$\mathbb{B}[0, \infty]$  is the category with a single object and the morphism set is  $[0, \infty]$  with addition as the composition.

is the unique lower semi-continuous convex functor that vanishes on the subcategory FP of optimal hypotheses (up to a non-negative constant).

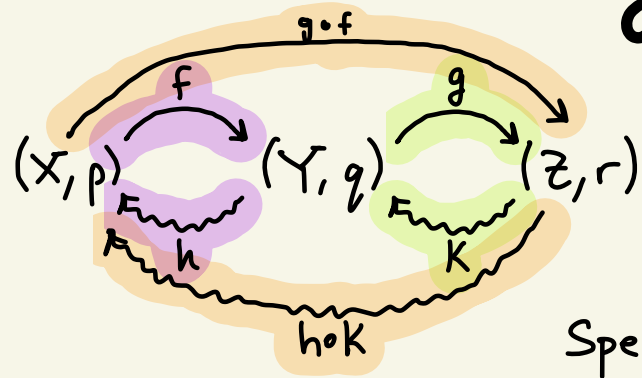
This is the main theorem of Baez and Fritz.

# Functoriality of Relative Entropy



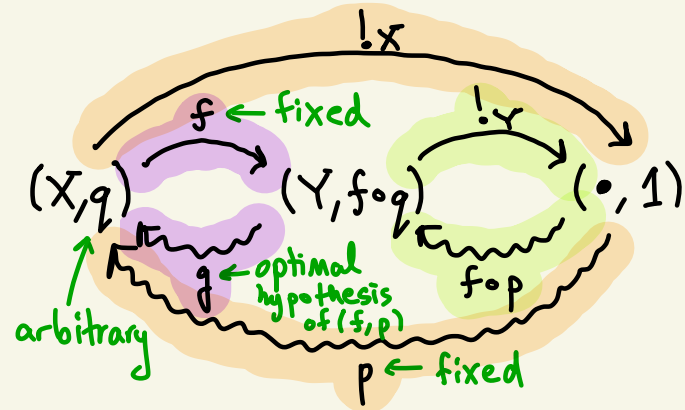
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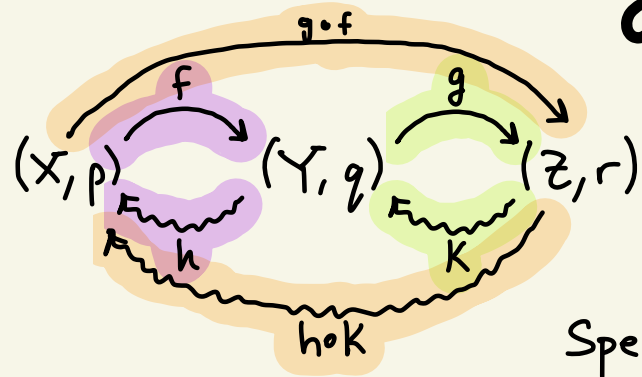


$$S(p \parallel h \circ g) + S(q \parallel K \circ r) = S(p \parallel h \circ K \circ r)$$

Special case of functoriality 1:

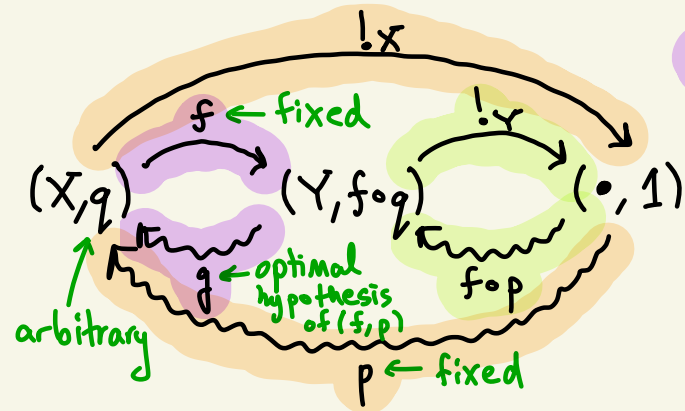


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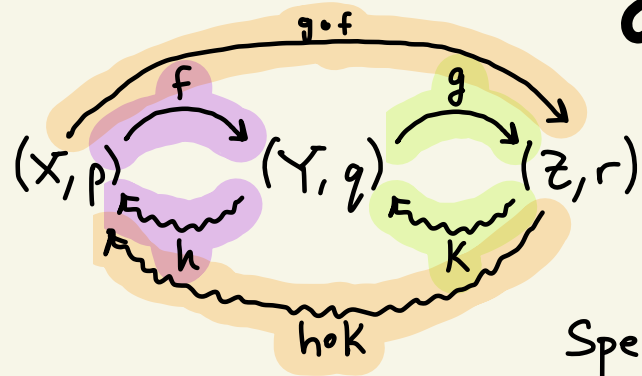
Special case of functoriality 1:



$$S(q \parallel g \circ f \circ g) + S(f \circ q \parallel f \circ p) = S(q \parallel p) \text{ i.e.,}$$

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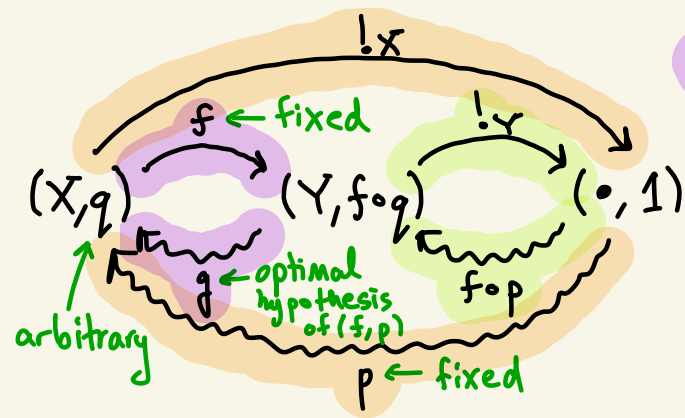
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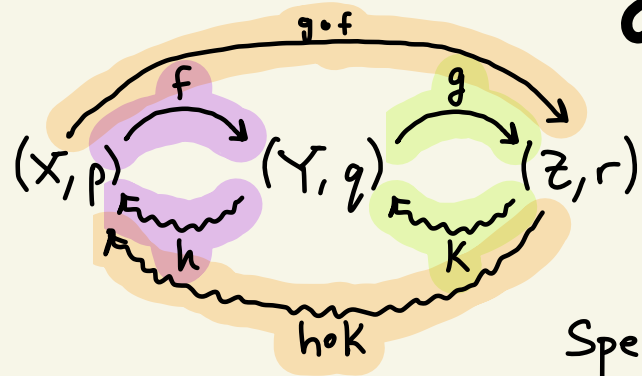
The conditional expectation property

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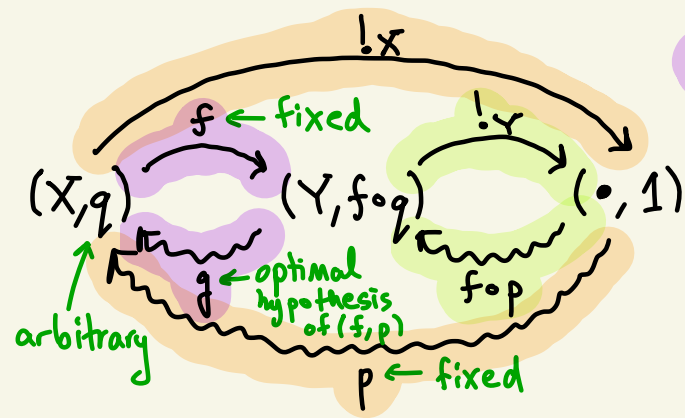
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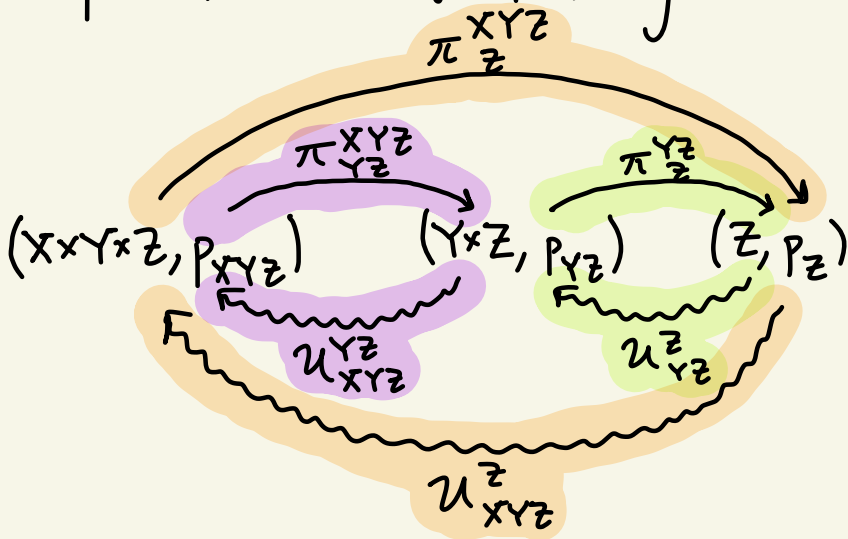
This term appears in the data-processing inequality (DPI). Since the RHS  $\geq 0$ , this is an improvement of the DPI.





# Functoriality of Relative Entropy

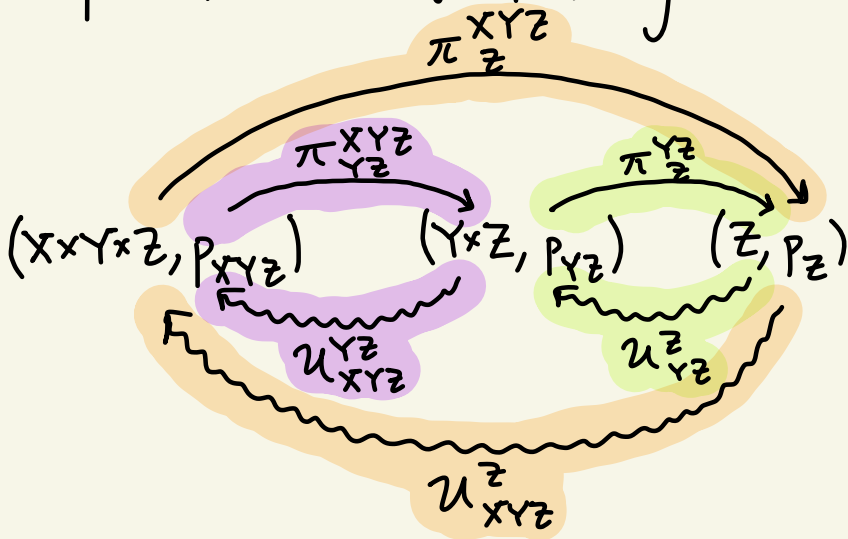
Special case of functoriality 2:



The  $\pi$ 's are projections, while the  $\nu$ 's are fiberwise uniform distributions.

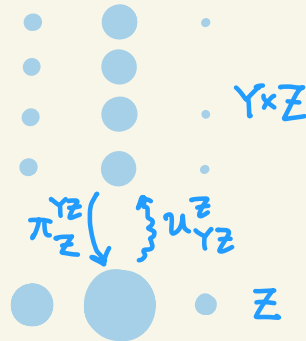
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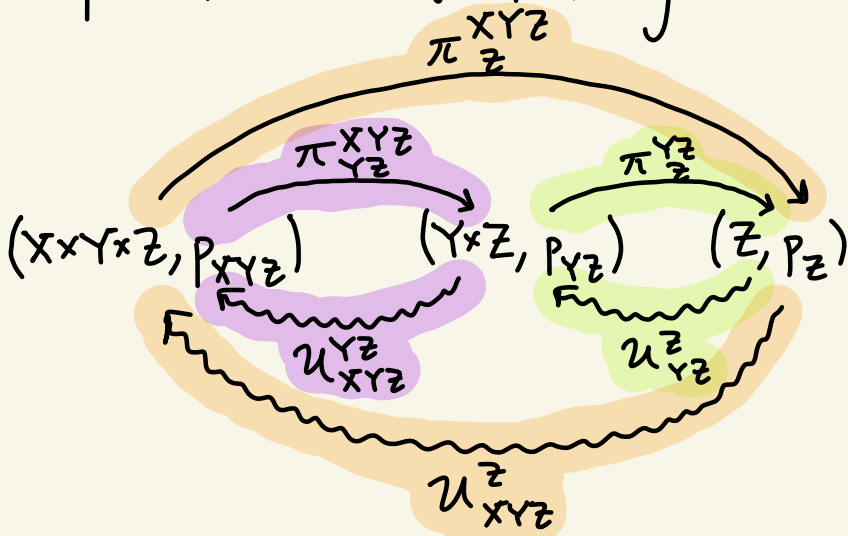
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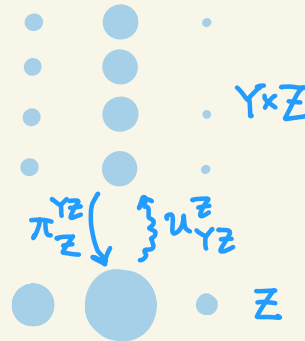
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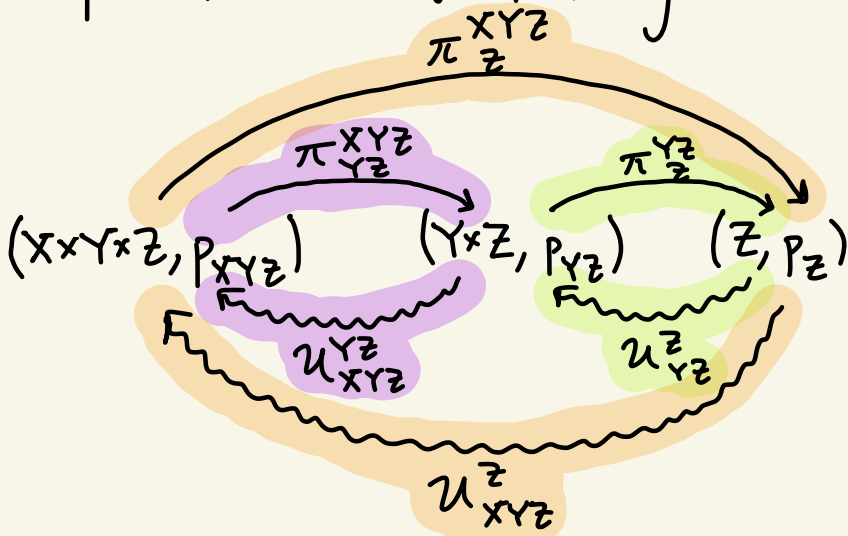
For example



$$S(P_{X Y Z} \parallel u_{X Y Z}^{Y Z} \circ P_{Y Z}) + S(P_{Y Z} \parallel u_{Y Z}^Z \circ P_Z) = S(P_{X Y Z} \parallel u_Z^{X Y Z} \circ P_Z)$$

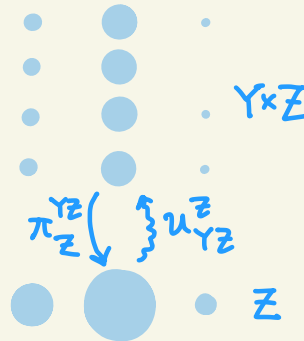
# Functoriality of Relative Entropy

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$$S(P_{XYZ} \parallel U_{XYZ}^{YZ} \circ P_{YZ}) + S(P_{YZ} \parallel U_{YZ}^Z \circ P_Z) = S(P_{XYZ} \parallel U_Z^{XYZ} \circ P_Z)$$

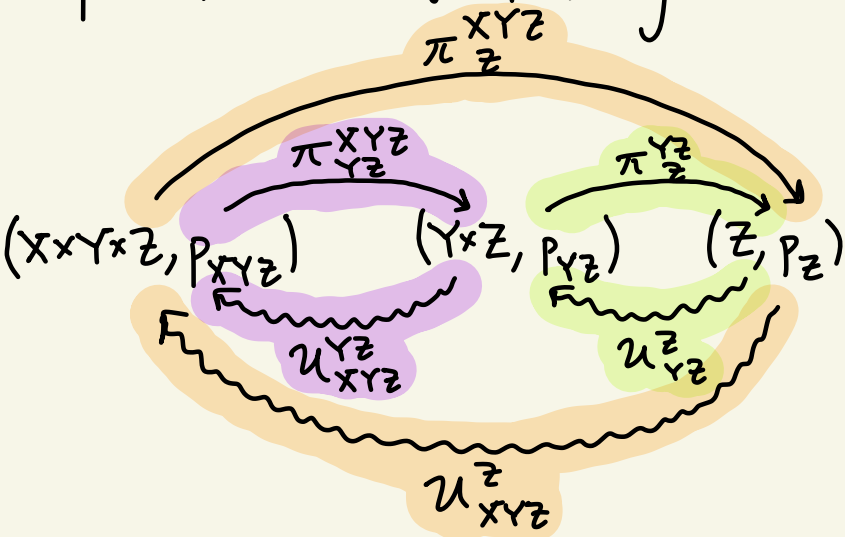
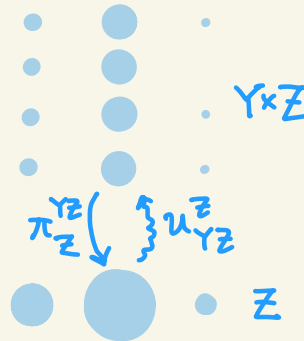
$$-H(X|YZ) + \log|X| - H(Y|Z) + \log|Y| = -H(XY|Z) + \log|X \times Y|$$

# Functoriality of Relative Entropy

Special case of functoriality 2: chain rule for conditional entropy

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$$H(X|YZ) + H(Y|Z) = H(XY|Z)$$

[Ref Cover-Thomas  
Eqn (2.21)]

# A stronger DPI

In our first example of functoriality, we found  
"given  $X \xrightarrow{f} Y$  and  $\bullet \rightsquigarrow X$ , set  $X \leftarrow \rightsquigarrow Y$  to be the  
optimal hypothesis, then  $S(q \| p) - S(f \circ q \| f \circ p) = S(q \| g \circ f \circ q) \quad \forall q$ ."

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Li and Winter proved (in 2014!)

"given  $X \xrightarrow{f} Y$  and  $\bullet \rightsquigarrow_p X$ , set  $X \leftarrow_q^g Y$  to be the Bayesian  
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Note that this implies monotonicity of relative entropy

$$S(q \| p) - S(f \circ q \| f \circ p) \geq 0,$$

which is sometimes called the data-processing inequality (DPI).

# Bayesian inverses

what's this?

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$$\begin{array}{ccc}
 X \xleftarrow{p} \bullet \xrightarrow{f \circ p} Y & & \\
 \Delta_X \downarrow & \equiv & \downarrow \Delta_Y \\
 X \times X \xrightarrow{\text{id}_X \times f} X \times Y \xleftarrow{g \times \text{id}_Y} Y \times Y & & 
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$$\begin{array}{ccc}
 X \xleftarrow{p} \bullet \xrightarrow{f \circ p} Y & \text{i.e.,} & f_{y \times p_x} = g_{x \times y}(f \circ p)_y \quad \forall x \in X, y \in Y \\
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i.e.,  $f_{y \times} p_x = g_{x \times y}(f \circ p)_y \quad \forall x \in X, y \in Y$

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# Bayesian inverses

what's this? ↘

"given  $X \xrightarrow{f} Y$  and  $\bullet \xrightarrow{p} X$ , set  $X \xleftarrow{g} Y$  to be the Bayesian inverse of  $(f, p)$ , then  $S(q \| p) - S(f \circ q \| f \circ p) \geq S(q \| g \circ f \circ q) \quad \forall q$ ."

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  - If  $g$  is an optimal hypothesis, then  $g$  is a Bayesian inverse and  $f$  is necessarily deterministic.



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- adjust Hawking's calculation to prove information conservation in black hole evaporation (formerly "the information paradox")



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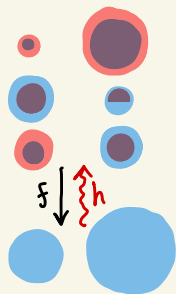
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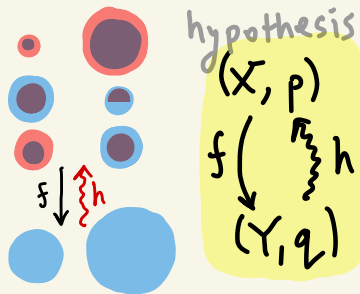
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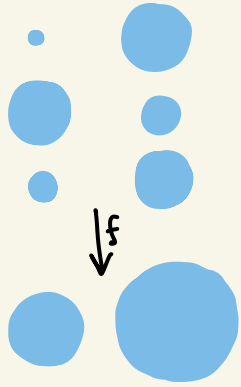
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# Subtle differences in quantum



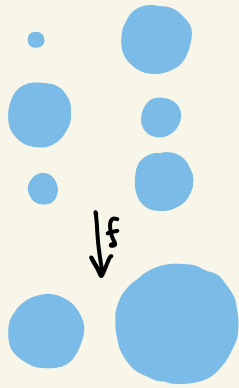
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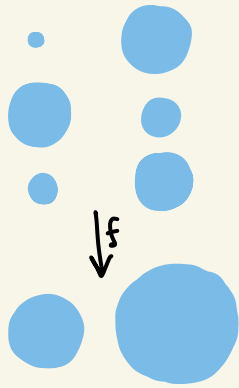
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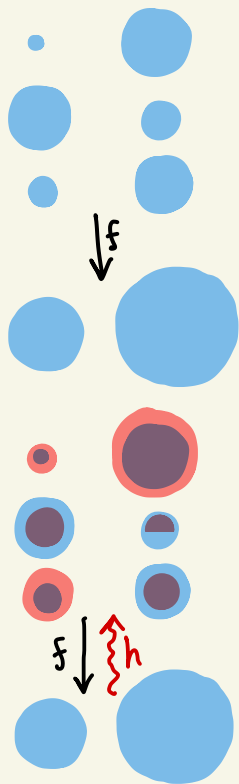
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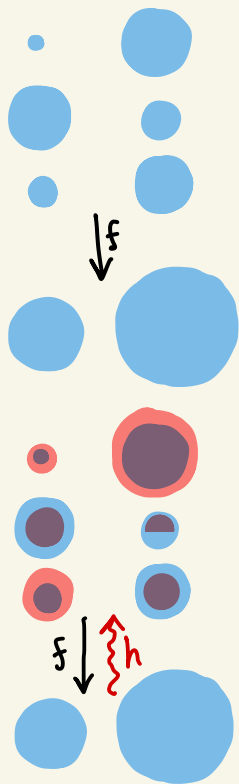
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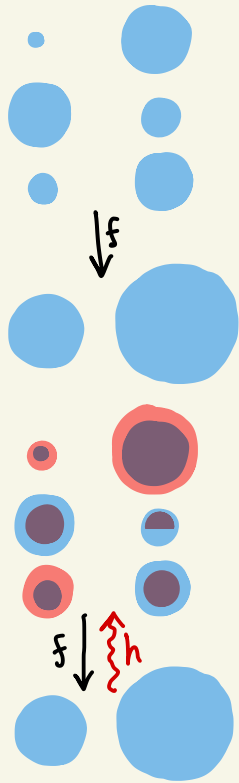
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Fact: Given  $(B, \xi) \xrightarrow{F} (A, \omega)$ , if an optimal

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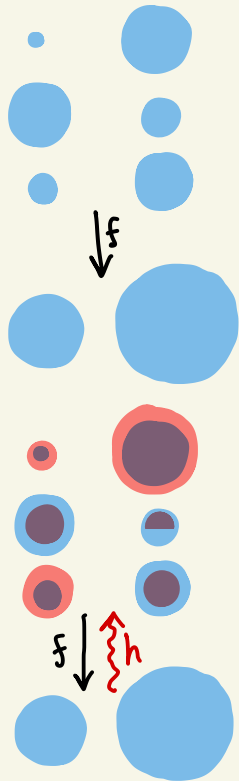
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See my talk @  
N. Yanofsky's seminar

This was one of the key observations in extending Baez-Fritz-Leinster's Theorem to quantum.

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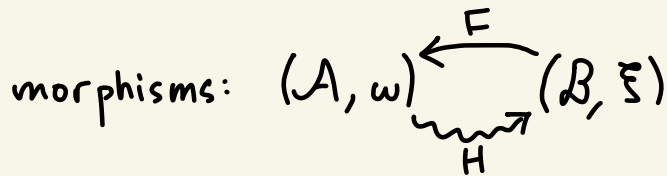
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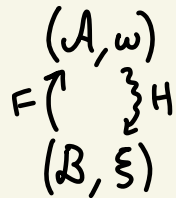


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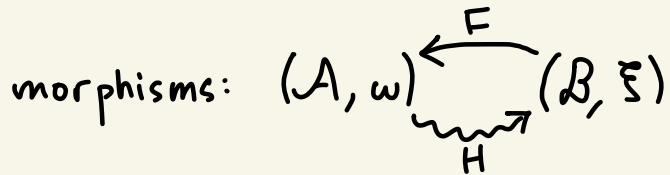
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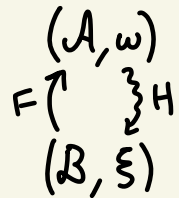


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In progress:

- 1) functoriality for all states
- 2) lower semi-continuity
- 3) characterization

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is false. Variants and weaker versions have been found.





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What connects all of these different approaches?

# Thank you!

## References

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