

Operadic Composition of Thermodynamical Systems

Owen Lynch

Universiteit Utrecht

May 13, 2022

Overview

1. What is entropy?
2. Thermo static systems
3. Composing thermostatic systems
(the central dogma)
4. ... with operads!

Examples

Conclusion

References

Part 1: What is Entropy?

Part 1, slide 1/3

Can point to many things called entropy...,

- Shannon Entropy
- Thermodynamic (empirically measured) entropy
- von Neumann entropy
- Rényi entropy

Why do we care about Entropy? Part 1, slide 2/3

- Maximizing it subject to a constraint often gives us the answer to a question we care about.
- What prior should we choose?
- Where is the equilibrium?
- What kind of large deviation is likely?

What is entropy?

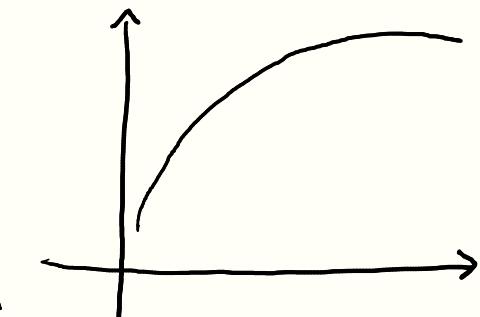
Part 1, slide 3/3

Without answering completely, we can make some assumptions.

- (1) Function of states
- (2) "Mixing" increases it.
- (3) Additive over independent subsystems
- (4) Something we want to maximize subject to constraints.

This talk will formalize the above picture

Formalize (1) + (2)



Def A thermostatic system consists of

- A convex space X
- A concave function $S: X \rightarrow \bar{\mathbb{R}} = \mathbb{R} \cup \{\pm\infty\}$

Ex $X = \{\text{prob distributions on } \{1 \dots n\}\}$

$$S(p) = - \sum_{i=1}^n p_i \log p_i$$

Why / what is convexity / concavity?

Def A **convex space** is a set X along with a function $c_\lambda : X \times X \rightarrow X$ for $\lambda \in [0,1]$ called the **mixing operation**, such that

- $c_1(x, y) = x$ (identity)
- $c_\lambda(x, x) = x$ (idempotency)
- $c_\lambda(x, y) = c_{1-\lambda}(y, x)$ (symmetry)
- $c_\lambda(c_\gamma(x, y), z) = c_{\lambda\gamma}(x, c_{1-\lambda}(y, z))$ (associativity)
if $\gamma(1-\lambda\gamma) = \lambda(1-\gamma)$

Ex Vector spaces, with $c_\lambda(x, y) = \lambda x + (1-\lambda)y$

Note Originally "Barycentric algebras" by Stone in 1949

Def/Ex A convex subspace $U \subseteq X$ of a convex space X is a subset $U \subseteq X$ such that $x, y \in U \Rightarrow c_\lambda(x, y) \in U \quad \forall \lambda \in [0, 1]$

Convex subspaces are convex spaces

Ex $\Delta^n = \{ p \in \mathbb{R}^{n+1} \mid \sum_{i=0}^n p_i = 1, p_i \geq 0 \forall i \}$

Examples of Convex Spaces

Part 2, slide 4/8

Ex If (A, \vee) is a semilattice, i.e.

- $a \vee a = a$
- $a \vee b = b \vee a$
- $(a \vee b) \vee c = a \vee (b \vee c)$

Then

$$c_\lambda(a, b) = \begin{cases} a & \text{if } \lambda = 1 \\ b & \text{if } \lambda = 0 \\ a \vee b & \text{otherwise} \end{cases}$$

makes A into a convex space.

Def Let $\bar{\mathbb{R}} = \mathbb{R} \cup \{\pm\infty\}$ have convex structure

$$c_\lambda(x, y) = (1-\lambda)x + \lambda y \quad \text{for } x, y \in \bar{\mathbb{R}}$$

$$c_\lambda(x, +\infty) = +\infty \quad \text{for } x \in \mathbb{R}, \lambda \in (0, 1]$$

$$c_\lambda(x, -\infty) = -\infty \quad \text{for } x \in \mathbb{R} \cup \{+\infty\}, \lambda \in (0, 1]$$

Note that $-\infty$ "beats" $+\infty$

Prop Every subset $U \subseteq \bar{\mathbb{R}}$ has a least upper bound.

Proof sup of an unbounded above set is $+\infty$

sup of an empty set is $-\infty$

Def A convex-linear function $f: X \rightarrow Y$ satisfies

$$f(c_\lambda(x, y)) = c_\lambda(f(x), f(y))$$

Def A convex relation is a convex subspace $R \subseteq X \times Y$

Ex $\{(x, f(x)) \mid x \in X\}$ is a convex relation iff f is convex-linear

Def Let Conv be the category of convex spaces and convex-linear functions, and let ConvRel be the category of convex spaces and convex relations

Def A concave function $s: X \rightarrow \bar{\mathbb{R}}$ satisfies

$$s(c_\lambda(x, y)) \geq c_\lambda(s(x), s(y))$$

Thermostatic systems revisited

Part 2 slide 7/8

Def A thermostatic system consists of

- Convex space X
- Concave function $S: X \rightarrow \bar{\mathbb{R}}$

Ex Gas has $X = \mathbb{R}_{>0}^3$, with coordinates (u, v, N) .
Entropy $S(u, v, N)$ encodes physical properties

$$\overset{\text{temperature}}{\underset{\text{temperature}}{\frac{1}{T}}} = \frac{\partial S}{\partial u} \quad \overset{\text{pressure}}{\underset{\text{pressure}}{\frac{\partial P}{T}}} = \frac{\partial S}{\partial v} \quad \overset{\text{chemical potential}}{\underset{\text{chemical potential}}{\frac{\partial \mu}{T}}} = \frac{\partial S}{\partial N}$$

Ex $X = \Delta^n$, Shannon entropy

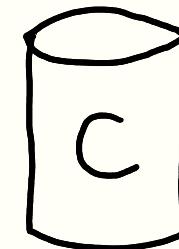
Ex $X = St(H)$, von Neumann entropy

More Examples of Thermostatic Systems Part 2 slide 8/8

Ex A water tank with heat capacity C has $X = \mathbb{R}_{>0}$, with coordinate U

$$S(U) = C \log U$$

$$\frac{1}{T} \stackrel{\text{def}}{=} \frac{\partial S}{\partial U} = \frac{C}{U} \Rightarrow U = CT$$



Note For large C , $U = CT$, small flux ΔU

Heat bath is "infinite heat capacity tank"

Ex A heat bath at temperature T has $X = \mathbb{R}$

$$S(\Delta U) = \frac{1}{T} \Delta U$$

Very simple, but not usually formalized!

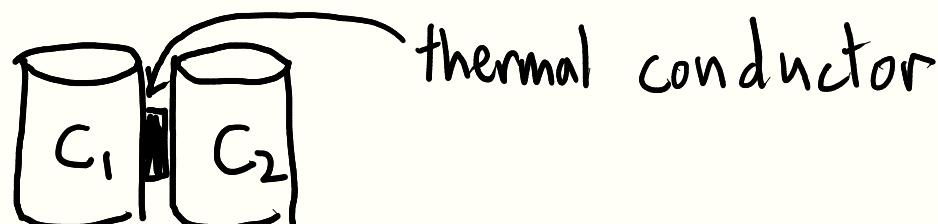
Part 3: Composition of Thermostatic Systems Part 3 slide 1/4

Recall assumptions about Entropy

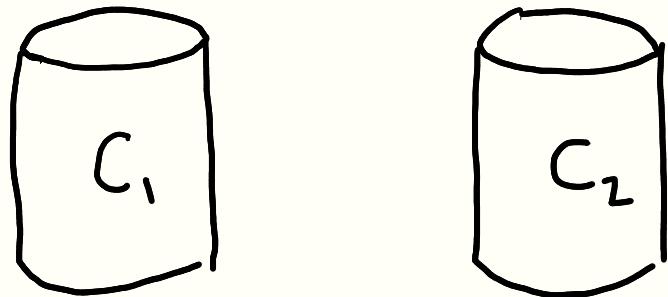
- (1) Function of states
- (2) Mixing increases it.
- (3) Additive over independent subsystems
- (4) Something we want to maximize subject to constraints.

We've formalized (1)+(2), now on to (3)+(4)

To start, we do an example: connecting two tanks



Apply (3) : Independent composition Part 3 slide 2/4



Start with $X_1 = R_{>0}$, $X_2 = R_{>0}$

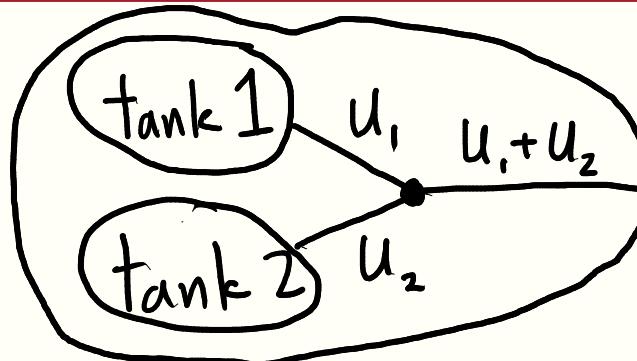
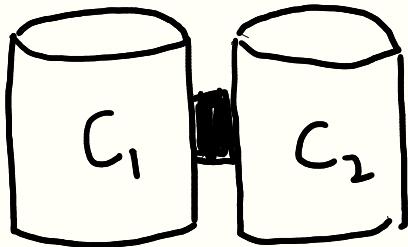
$$S_1(u_1) = C_1 \log(u_1)$$

$$S_2(u_2) = C_2 \log(u_2)$$

Make $X = X_1 \times X_2$,

$$S(u_1, u_2) = S_1(u_1) + S_2(u_2) = C_1 \log u_1 + C_2 \log u_2$$

Apply (4): Constrain with Conserved Quantities Part 3, slide 3/4



- Total energy $u = u_1 + u_2$ is conserved.
- Relation $R \subseteq (R_{>0} \times R_{>0}) \times R_{>0}$
 $((u_1, u_2), u) \in R \Leftrightarrow u_1 + u_2 = u$
- Entropy of u is supremum of entropies of compatible states

$$S(u) = \sup_{u_1 + u_2 = u} S_1(u_1) + S_2(u_2)$$
- At supremum, $\frac{\partial S_1}{\partial u_1} = \frac{\partial S_2}{\partial u_2}$, i.e. $T_1 = T_2$
- New system $(R_{>0}, S)$, $S(u) = (C_1 + C_2) \log(u) + K$

Central Dogma of Thermodynamics Part 3, slide 4/4

Independent composition:

$$(x_1, s_1), \dots, (x_n, s_n) \mapsto (x_1 \times \dots \times x_n, s_1 + \dots + s_n)$$

Maximize over constraints

$$(x, s), R \subseteq X \times Y \xrightarrow{\text{convex}} (Y, R_* s)$$

$$R_* s(y) = \sup_{\substack{(x,y) \in R \\ \text{concave!}}} s(x)$$

Composition =

Independent composition

↓
Maximize over constraints

Aside

Central Dogma bridges Thermo + Stat Mech

Ω finite, $X = P(\Omega)$, $S: X \rightarrow \bar{\mathbb{R}}$ Shannon

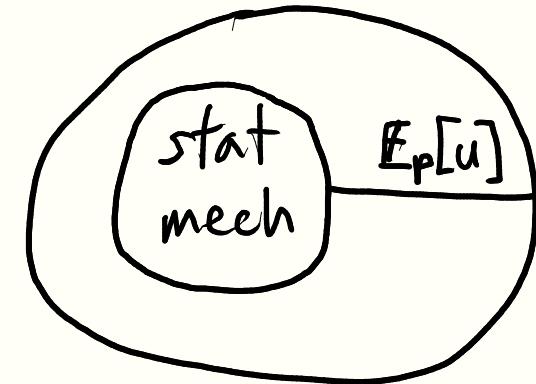
(X, S) is "stat mech-esque"

$H: \Omega \rightarrow \mathbb{R}_{>0}$, $Y = \mathbb{R}_{>0}$

$R \subseteq X \times Y$, $(p, u) \in R \Leftrightarrow E_p[H] = u$

$$R_*S(u) = \sup_{E_p[H] = u} S(p)$$

(Y, R_*S) is "thermo-esque"



supremum
is canonical
distribution!

Raw materials:

Def ConvRel

- convex spaces as objects
- convex relations $R \subseteq X \times Y$ as morphisms

Def Ent: ConvRel \rightarrow Set

$$\text{Ent}(X) = \{S: X \rightarrow \bar{\mathbb{R}}, S \text{ concave}\}$$

$$\text{Ent}(R) = S \mapsto R_* S \quad (R_* S(y) = \sup_{(x,y) \in R} S(x))$$

Def Natural transformation

$$\text{ind}_{X_1, \dots, X_n}: \text{Ent}(X_1) \times \dots \times \text{Ent}(X_n) \rightarrow \text{Ent}(X_1 \times \dots \times X_n)$$

$$(S_1, \dots, S_n) \mapsto S_1 + \dots + S_n$$

$$\text{Note: } \text{ind}_\phi: 1 \rightarrow \text{Ent}(\mathbb{R}^0) \quad * \mapsto 0$$

Def A symmetric monoidal category consists of

- A category C
- A functor $\otimes : C \times C \rightarrow C$
- An object $I \in C_0$

such that (assoc, unit, sym, etc.)

Ex Any category with finite (co)products

Ex $(\text{Conv}, \times, R^0)$

Ex $(\text{ConvRel}, \times, R^0)$ (because $\text{Conv} \overset{\text{wide}}{\subseteq} \text{ConvRel}$)

Note: \times NOT product in ConvRel

Def A lax symmetric monoidal functor from a SMC C to a SMC D is..

- A functor $F: C \rightarrow D$
- A natural transformation

$$\epsilon_{x_1, \dots, x_n}: F(x_1) \otimes_D \cdots \otimes_D F(x_n) \rightarrow F(x_1 \otimes_C \cdots \otimes_C x_n)$$

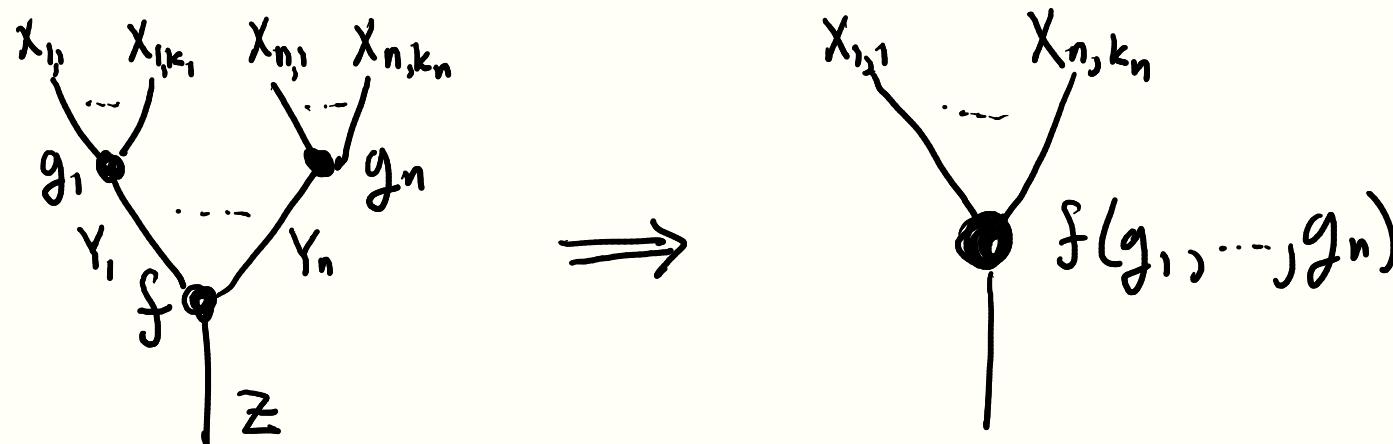
for $n \in \mathbb{N}$ (when $n=0$, $\epsilon_\emptyset: I_D \rightarrow F(I_C)$)
..., satisfying coherence conditions...

Prop (Ent, ind) is a lax symmetric monoidal functor
 $(\text{Conv Rel}, \times, R^\circ) \rightarrow (\text{Set}, \times, 1)$

Recall A (multisorted) operad \mathcal{O} has

- a set of **types** \mathcal{O}_0
- for every $X_1, \dots, X_n, Y \in \mathcal{O}_0$,
a set of **operations** $\mathcal{O}(X_1, \dots, X_n; Y)$

With (suitably well-behaved) composition



Def Given a SMC $(\mathcal{C}, \otimes, \mathbb{I})$ define $\text{Op}(\mathcal{C})$

$$\text{Op}(\mathcal{C})_0 = \mathcal{C}_0$$

$$\text{Op}(\mathcal{C})(x_1, \dots, x_n; Y) = \text{Hom}_{\mathcal{C}}(x_1 \otimes \dots \otimes x_n, Y)$$

Ex Let $\mathcal{CR} = \text{Op}(\text{ConvRel})$

- types are convex sets
- operations are convex relations

$$R \subseteq (X_1 \times \dots \times X_n) \times Y$$

Description of how to compose $(x_1, s_1) \dots (x_n, s_n)$!

Def An operad algebra F of an operad \mathcal{O} is

- A set $F(x)$ for every type $x \in \mathcal{O}_0$
- A function $F(f) : F(x_1) \times \cdots \times F(x_n) \rightarrow F(Y)$

for every operation $f \in \mathcal{O}(x_1, \dots, x_n; Y)$

respecting composition, identities, symmetries...

Prop SMC $(\mathcal{C}, \otimes, I)$, LSM functor $F : \mathcal{C} \rightarrow \text{Set}$,
 then $Op(F)$ is operad algebra of $Op(\mathcal{C})$

$$Op(F)(X) = F(X) \quad X \in Op(\mathcal{C})_0 = \mathcal{C}_0$$

$$Op(F)(f) = F(x_1) \otimes \cdots \otimes F(x_n) \xrightarrow{\epsilon_{x_1, \dots, x_n}} F(x_1 \otimes \cdots \otimes x_n) \xrightarrow{F(f)} F(Y)$$

$$f \in Op(\mathcal{C})(x_1, \dots, x_n; Y) = \text{Hom}_{\mathcal{C}}(x_1 \otimes \cdots \otimes x_n, Y)$$

Thm $\text{Op}(\text{Ent})$ is an operad algebra
of $\text{Op}(\text{ConvRel})$.

The action of a relation $R \subseteq X_1 \times \dots \times X_n \times Y$
on entropy functions S_1, \dots, S_n is

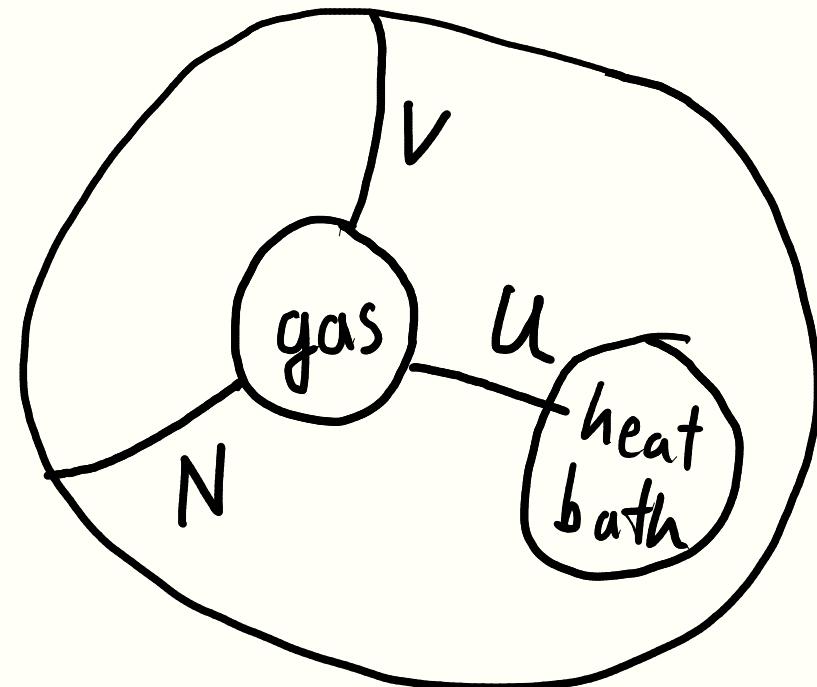
$$\begin{aligned} \text{Ent}(x_1) \times \dots \times \text{Ent}(x_n) &\xrightarrow{\text{ind}} \text{Ent}(x_1 \times \dots \times x_n) \xrightarrow{R^*} \text{Ent}(Y) \\ (S_1, \dots, S_n) &\longmapsto S_1 + \dots + S_n \longmapsto R^*(S_1 + \dots + S_n) \end{aligned}$$

- The central dogma is precisely
the rule for making an operad algebra
out of a LSM functor!

Example 1/3 : Ideal gas at constant T

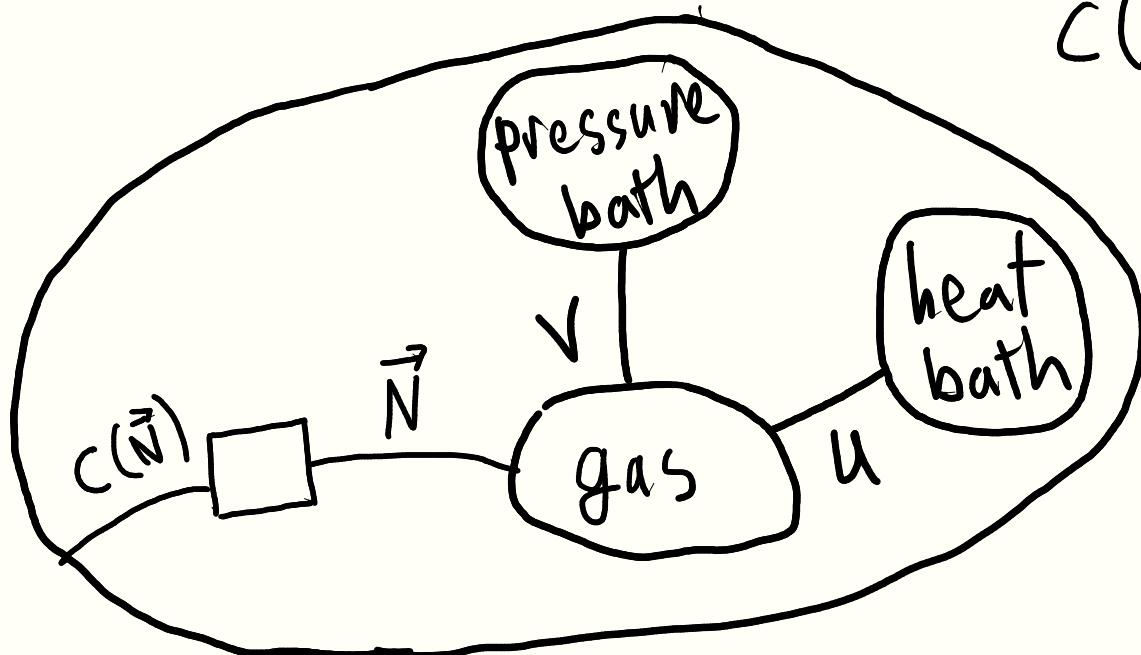
gas
 $X_1 = \mathbb{R}_{>0}^3$, $X_2 = \mathbb{R}$, $Y = \mathbb{R}_{>0}^2$

heat bath
gas at const. T



supremum happens at $\frac{\partial S}{\partial u} = \frac{1}{T}$

Example 2/3: Chemical Reactions



$C(\vec{N}) = \underline{\text{conserved quantities}}$



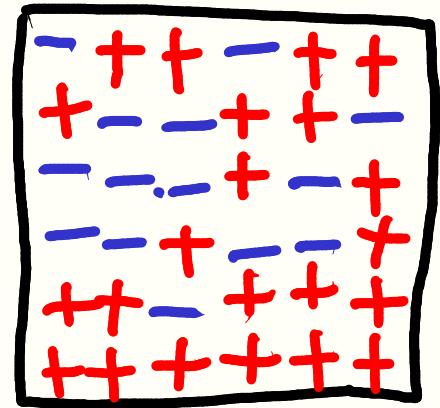
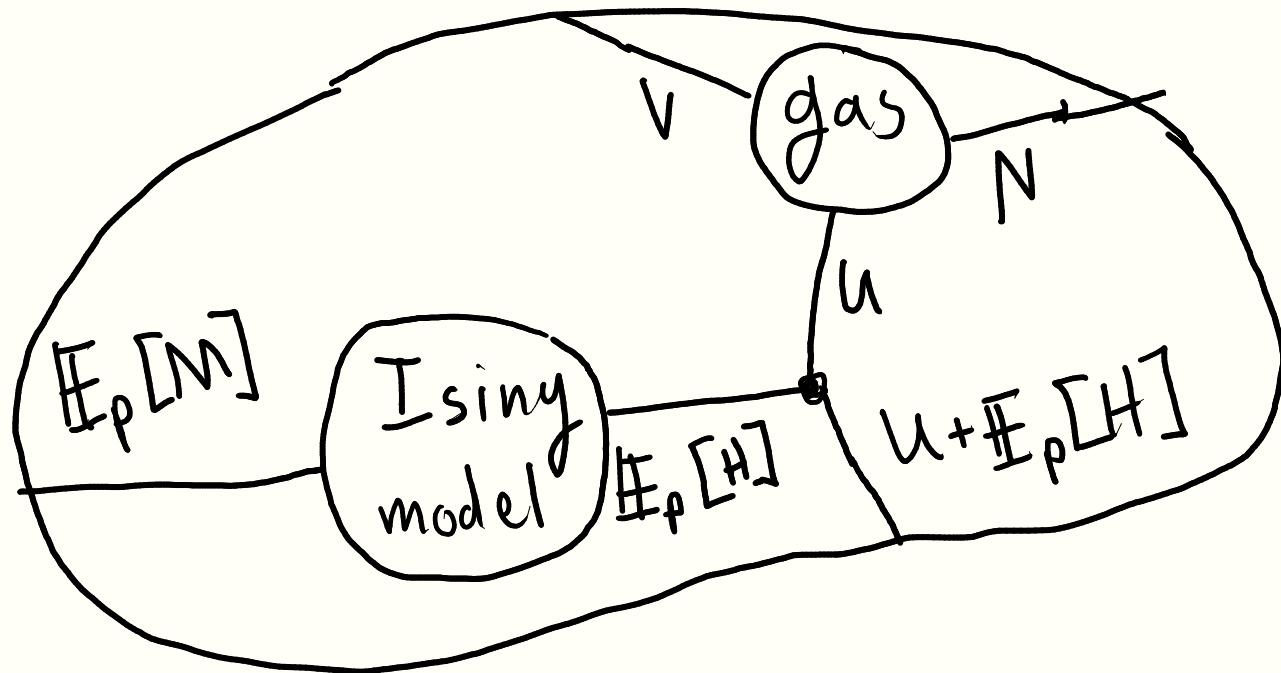
$[E]$ is conserved

$[S] + [P]$ is conserved

Maximum entropy tells us U, V, \vec{N}

given $T, P, [E], [S] + [P]$

Example 3/3: Ising Model + Gas



- Traditional statistical mechanics uses a heat bath
- We can replace the heat bath with more complex systems, like gas

Conclusions

- We can do a lot with our assumptions!
- Having different entropies in same framing allows composition of previously incompatible systems.
- "central dogma" = build an operad algebra from a lax symmetric monoidal functor
- Other physical systems can be composed in this way

References

Compositional Thermostatics,
John Baez, Owen Lynch, and Joe Moeller
(find it on arXiv)